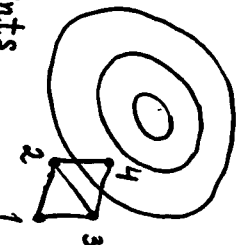


Algorithms for unconstrained optimization on \mathbb{R}^n $F = \mathbb{R}^n$ — no constraintsSimplex algorithmA simplex in \mathbb{R}^n = set of $n+1$ equidistant pointse.g. in \mathbb{R}^2 :

Basic simplex alg. — not the same as the Simplex Alg. of Linear Programming

- start with simplex (1, 2, 3)
- reflect the point in the simplex (1, 2, 3) which has greatest cost in the centroid of the remaining points giving a new simplex (2, 3, 4) "nearer" a l.m.

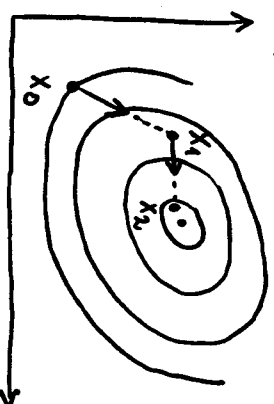


- repeat for simplex (2, 3, 4) etc.

Properties: simple in principle but needs modifications to help convergence to a l.m.

Search direction algorithm

$$x_1 = x_0 + w_0 s_0$$



- Choose $x_0 \in \mathbb{R}^n$ (an estimate of a g.m. \hat{x})
- Set iteration index $j=0$
- 1) [iteration j]:

- Using info available at x_j , choose a search direction $s_j \in \mathbb{R}^n$
- choose a scalar $w_j > 0$ so $V(x_j + w_j s_j) < V(x_j)$
- Set $x_{j+1} := x_j + w_j s_j$
 $j := j+1$ and go to 1)

choose s_j so V should decrease as one moves along s_j , away from x_j

Remarks:

very important type of alg. as it can be developed & adapted a lot

Problems:

choice of s_j, w_j

Concerning choice of s_j :

The Gradient

\mathbb{R}^2 case

Under weak conditions \leftarrow

see cond. (a) & (b) on p. 2.4

$$V\left[\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} + \begin{pmatrix} \delta x_1 \\ \delta x_2 \end{pmatrix}\right] \cong V\left[\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}\right] + \frac{\partial V}{\partial x_1}[x] \delta x_1 + \frac{\partial V}{\partial x_2}[x] \delta x_2$$

$$= V[x] + \left[\frac{\partial V[x]}{\partial x_1} \quad \frac{\partial V[x]}{\partial x_2} \right]^T \begin{pmatrix} \delta x_1 \\ \delta x_2 \end{pmatrix}$$

$\underbrace{\hspace{10em}}_{\nabla V[x]} \quad \underbrace{\hspace{2em}}_{\delta x}$

i.e.

the gradient of V at x

$$V[x + \delta x] \cong V[x] + \nabla V[x]^T \delta x$$

for small δx

For \mathbb{R}^n case:

(2.1) Th Consider $V: \mathbb{R}^n \rightarrow \mathbb{R}$ when

the p.ds. $\frac{\partial V}{\partial x_i}[z]$:

(a) all exist $\forall z$ s.t. $\|z - x\| < \epsilon$ #
for some $\epsilon > 0$

(b) all are continuous at $z = x$

Then:

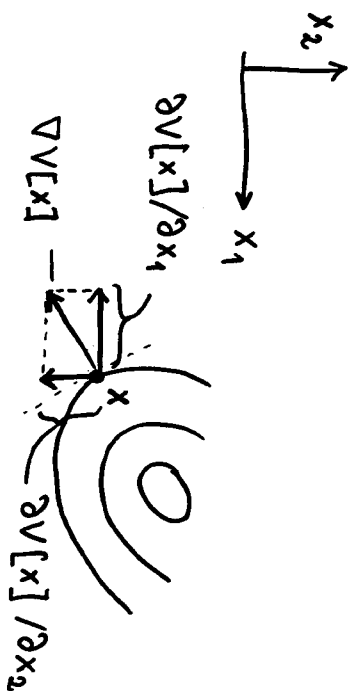
V is called differentiable at x and, for small δx :

$$V[x + \delta x] \cong V[x] + \nabla V[x]^T \delta x$$

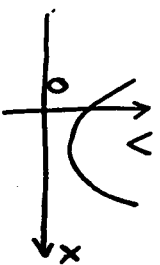
$$\nabla V[x] \triangleq \begin{bmatrix} \frac{\partial V[x]}{\partial x_1} \\ \vdots \\ \frac{\partial V[x]}{\partial x_n} \end{bmatrix} \in \mathbb{R}^n$$

the gradient of V at x

i.e. the p.ds. all exist everywhere near z

Graphical representation of ∇V (2.2) Fact :Even for the \mathbb{R}^n case :

$\nabla V[x]$ is orthogonal to the constant V contour through x and points uphill

Quadratics on \mathbb{R}^n :Quadratic of $x \in \mathbb{R}$ 

$$V(x) = a + bx + \frac{1}{2}x^T C x$$

$\begin{matrix} \mathbb{R} & \mathbb{R} & \mathbb{R} \\ a & b & C \end{matrix}$

$\left[\begin{array}{l} C \in \mathbb{R}^{n \times n} \text{ \& symmetric} \\ \text{for a min : } C > 0 \end{array} \right] \#$

in the sense that a scalar c can be thought of as a $[x]$ matrix, which is symmetric

Standard quadratic of $x \in \mathbb{R}^n$:

$$V(x) = a + b^T x + \frac{1}{2} x^T C x$$

$\begin{matrix} \mathbb{R} & \mathbb{R}^n & \mathbb{R}^n & \mathbb{R}^n \\ a & b & C \end{matrix}$

$C \in \mathbb{R}^{n \times n}$

read as :
 C is positive
 definite & symmetric

Tests for $C^T = C > 0$ For symmetric C :

$$[C > 0] \Leftrightarrow \left\{ \begin{array}{l} C_{11} > 0 ; \\ \begin{vmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{vmatrix} > 0 ; \\ \begin{vmatrix} C_{11} & C_{12} & C_{13} \\ C_{21} & C_{22} & C_{23} \\ C_{31} & C_{32} & C_{33} \end{vmatrix} > 0 ; \\ \dots \\ |C| > 0 \end{array} \right.$$

i.e. $x^T C x > 0, \forall x \neq 0$
 \mathbb{R}^n

$$C^T = C \text{ \& } C > 0$$

Also:

$$[C > 0] \Leftrightarrow \lambda_i[C] > 0, \forall i$$

the eigenvalues of C ;
since $C^T = C$, all the eigenvalues of C are real.
This says that for C to be positive definite, all its eigenvalues must be strictly positive

By simple algebra:

$$V[x + \delta x] = V[x] + [b + Cx]^T \delta x + \frac{1}{2} \delta x^T C \delta x \quad (2.4)$$

$$\stackrel{\text{second order}}{\approx} V[x] + [b + Cx]^T \delta x$$

for small δx

Suggests #:

$$\nabla V[x] = b + Cx$$

check by partial differentiation

because we also have (from Th.(2.1)):

$$V(x + \delta x) \approx V(x) + \nabla V(x)^T \delta x$$

Hence:

(2.5) Th: For a standard quadratic on \mathbb{R}^n

- (i) V is differentiable on \mathbb{R}^n
(i.e. V is differentiable at x , $\forall x \in \mathbb{R}^n$)
- (ii) $\nabla V[x] = b + Cx$
- (iii) $V[x + \delta x] = V[x] + \nabla V[x]^T \delta x + \frac{1}{2} \delta x^T C \delta x$

An optimality condition for general V (2.6) Th If $V: \mathbb{R}^n \rightarrow \mathbb{R}$ is differentiable at \hat{x} :

$$[\hat{x} = \text{a l.m. of } V \text{ on } \mathbb{R}^n] \Rightarrow [\nabla V(\hat{x}) = 0]$$

unconstrained

necessary condition for $\hat{x} = \text{a l.m.}$ Remark:

So, one way to find a l.m. is:

solve $\nabla V(\hat{x}) = 0$ for \hat{x} check if $\hat{x} = \text{a l.m.}$

Proof of Th. (2.6):contrapositive i.e.:

$$\neg [\nabla V(\hat{x}) = 0] \Rightarrow \neg [\hat{x} = a \text{ l.m.}] :$$

$$\neg [\nabla V(\hat{x}) = 0] \Rightarrow [\nabla V(\hat{x}) \neq 0]$$

 \Rightarrow for all sufficiently small real $\omega > 0$

$$(i) \quad \hat{x} + \omega [-\nabla V(\hat{x})] = \text{near } \hat{x}$$

$$(ii) \quad V(\hat{x} + \omega [-\nabla V(\hat{x})]) =$$

$$\stackrel{\textcircled{\#}}{\cong} V(\hat{x}) + \nabla V(\hat{x})^T \delta x$$

$$= V(\hat{x}) - \omega \|\nabla V(\hat{x})\|^2 < V(\hat{x})$$

$$\Rightarrow [\hat{x} \neq a \text{ l.m.}] \Rightarrow \neg [\hat{x} = a \text{ l.m.}]$$

QED

the main point here is that the fact that
 $V(\hat{x}) + \nabla V(\hat{x})^T \delta x < V(\hat{x})$
 guarantees that
 $V(\hat{x} + \delta x) < V(\hat{x})$

Application to minimization of a standard quadratic

$$\textcircled{I} \quad \text{Solve } \nabla V(\hat{x}) = 0$$

$$\parallel \begin{matrix} b + C\hat{x} \end{matrix}$$

Th. (2.5 (ii))

Fact : $(C^T = C > 0) \Rightarrow C^{-1}$ exists

Hence :

$$\hat{x} = -C^{-1}b$$

$$\textcircled{II} \quad \text{Check : } \hat{x} = a \text{ l.m. ?}$$

$$V(x) = V(\hat{x} + [x - \hat{x}]) \stackrel{\text{Th. 2.5 (ii)}}{=} \underbrace{V(\hat{x})}_{\delta x} + \nabla V(\hat{x})^T \delta x + \frac{1}{2} \delta x^T C \delta x$$

$$\stackrel{\parallel}{=} 0 \stackrel{\gg}{\geq} 0 \text{ as } C > 0$$

$$\geq V(\hat{x}), \forall x \in \mathbb{R}^n$$

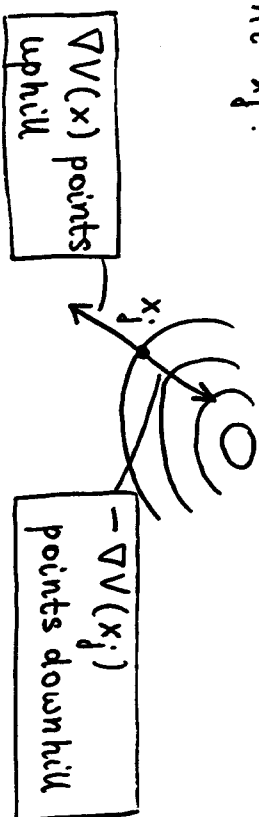
so $\hat{x} = a$ g.m. of V on \mathbb{R}^n \Downarrow

(2.7)

Th : $\hat{x} = -C^{-1}b$ is a unique g.m. of a quadratic V on \mathbb{R}^n and :

$$V(\hat{x}) = a - \frac{1}{2} b^T C^{-1} b$$

not just
a l.m.

More about the Search Directions Alg.A choice for s_j :Assume that V is differentiable on \mathbb{R}^n At x_j :choose: $s_j = -\nabla V(x_j)$ Steepest-descent
search direction

as then s_j points downhill
so that V should decrease as
we move along s_j from x_j

The steepest-descent alg.Choose $x_0 \in \mathbb{R}^n$. Set $j = 0$ 1) [At iteration j]

- Stop if $\nabla V(x_j) = 0$ -

in practice, use
 $\|\nabla V(x_j)\| < 10^{-\epsilon}$ say

as then x_j satisfies
a necessary condition
for a l.m.

- Set: $s_j := -\nabla V(x_j)$

- Choose a scalar $w_j > 0$ so

$$V(x_j + w_j s_j) < V(x_j)$$

often, w_j chosen to
minimize $V(x_j + w_j s_j)$
 $w \in \mathbb{R}$
- at least approximately

- Set: $x_{j+1} := x_j + w_j s_j$
 $j := j+1$

- Go to 1)

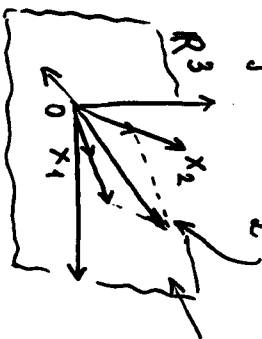
Some revision of linear algebra:

Consider $x^1, x^2, \dots, x^k \in \mathbb{R}^n$

A linear combination of these is $\sum_{i=1}^k d_i x^i$

$\underbrace{\quad}_{\text{scalars}}$

e.g. $2x^1 + \frac{1}{2}x^2$



plane through $0, x^1, x^2$
= set of all linear combinations
of x^1, x^2

$$= \{d_1 x^1 + d_2 x^2 : d_1, d_2 \in \mathbb{R}\}$$

the set of vectors $d_1 x^1 + d_2 x^2$ that one gets as d_1, d_2 vary all over the set of all real numbers

$$\mathcal{L}[x^1, x^2]$$

the linear subspace spanned by x^1, x^2

We can say: $x^1, \dots, x^k \in \mathbb{R}^n$ are linearly independent iff:

$$\{0\} \subsetneq \mathcal{L}[x^1] \subsetneq \mathcal{L}[x^1, x^2] \subsetneq \dots \subsetneq \mathcal{L}[x^1, \dots, x^k]$$

Then we say:

$\mathcal{L}[x^1]$ has dimension 1

$\mathcal{L}[x^1, x^2]$ " " 2

$\mathcal{L}[x^1, x^2, x^3]$ " " 3 etc.

Further: $x^1, \dots, x^n \in \mathbb{R}^n$ are linearly indep.

$$\text{iff } \mathcal{L}[x^1, \dots, x^n] = \mathbb{R}^n$$

In general:

$$\mathcal{L}[x^1, \dots, x^k]$$

the linear subspace spanned by x^1, \dots, x^k

$$\left\{ \sum_{i=1}^k d_i x^i : d_i \in \mathbb{R}, \forall i \right\} \subset \mathbb{R}^n$$

the plane through $0, x^1, \dots, x^k \subset \mathbb{R}^n$

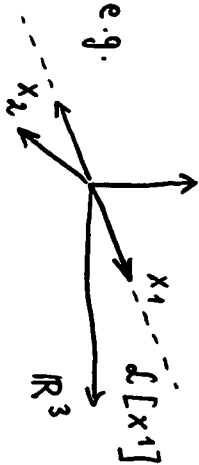
Further: $x^1, \dots, x^k \in \mathbb{R}^n$ are linearly indep.

$$\text{iff: } \sum_{i=1}^k d_i x^i = 0 \Rightarrow d_i = 0 \forall i$$

Sometimes

$$\mathcal{L}[x^1] = \mathcal{L}[x^1, x^2]$$

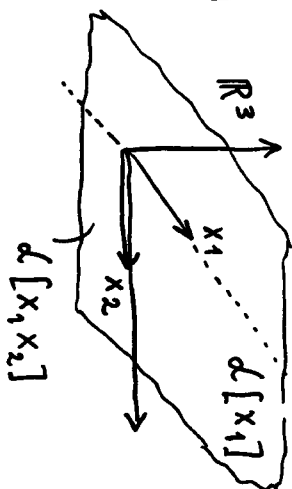
e.g.



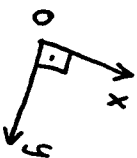
But usually

$$\mathcal{L}[x^1] \neq \mathcal{L}[x^1, x^2]$$

e.g.



Vectors $x, y \in \mathbb{R}^n$ are orthogonal



$$\text{iff } x^T y = 0$$