

# lecture 3

## Combinational logic 1

- truth tables
- Boolean algebra
- sum of products and product-of-sums
- logic gates

January 18, 2016

# Quiz 1

Class should start after ~15 min.

Let  $A, B$  be binary variables  
("boolean")

$1 \equiv \text{true}, \quad 0 \equiv \text{false}$

Notation:

$A \cdot B$	$\equiv$	$A$ and $B$
$A + B$	$\equiv$	$A$ or $B$
$\overline{A}$	$\equiv$	not $A$

One uses  $+, \cdot$  instead of  $\vee, \wedge$ .  
which you may have seen elsewhere.

# Truth Tables

Notation:  $A \cdot B \equiv A \text{ and } B$   
 $A + B \equiv A \text{ or } B$   
 $\overline{A} \equiv \text{not } A$

A	B	$A \cdot B$	$A + B$
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	1

A	$\overline{A}$
0	1
1	0

(exclusive or)

NAND

NOR

XOR

A	B	$\overline{A \cdot B}$	$\overline{A + B}$	$A \oplus B$
0	0	1	1	0
0	1	1	0	1
1	0	1	0	1
1	1	0	0	0

There are  $2^4 = 16$  possible boolean functions.

$$f: \{0,1\} \times \{0,1\} \rightarrow \{0,1\}$$

A	B	$\gamma_1$	$\gamma_2$	$\gamma_3$	...	$\gamma_{16}$
0	0					
0	1					
1	0					
1	1					

We typically only work with AND, OR, NAND, NOR, XOR.

# Laws of Boolean Algebra

identity

$$A + 0 = A$$

$$A \cdot 1 = A$$

inverse

$$A + \bar{A} = 1$$

$$A \cdot \bar{A} = 0$$

one and zero

$$A + 1 = 1$$

$$A \cdot 0 = 0$$

commutative

$$A + B = B + A$$

$$A \cdot B = B \cdot A$$

associative

$$(A + B) + C = A + (B + C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

distributive

$$A \cdot (B + C) = (A \cdot B) + (A \cdot C)$$

$$A + (B \cdot C) = (A + B) \cdot (A + C)$$

de Morgan

$$\overline{(A + B)} = \bar{A} \cdot \bar{B}$$

$$\overline{A \cdot B} = \bar{A} + \bar{B}$$

# Laws of Boolean Algebra

distributive

$$A + (B \cdot C) = (A + B) \cdot (A + C)$$

Note this one behaves differently from integers or reals.



# Example

$$Y = \overline{A \cdot B \cdot C} \cdot (A \cdot B + A \cdot C)$$

A	B	C	Y
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

$$Y = \overline{A \cdot B \cdot C} \cdot (A \cdot B + A \cdot C)$$

A	B	C	$A \cdot B \cdot C$	$\overline{A \cdot B \cdot C}$	$A \cdot B$	$A \cdot C$	$A \cdot B + A \cdot C$	Y
0	0	0	0	1	0	0	0	0
0	0	1	0	1	0	0	0	0
0	1	0	0	1	0	0	0	0
0	1	1	0	1	0	0	0	0
1	0	0	0	1	0	0	0	1
1	0	1	0	1	0	1	1	1
1	1	0	0	1	1	0	1	1
1	1	1	1	0	1	1	1	0



# Sum of Products

$$Y = \overline{A \cdot B \cdot C} \cdot (A \cdot B + A \cdot C)$$
$$= A \cdot \overline{B} \cdot C + A \cdot B \cdot \overline{C}$$

Q: For 3 variables A, B, C, how many terms can we have in a sum of products representation ?

A:  $2^3 = 8$  i.e. previous slide

$$Y = A \cdot \bar{B} \cdot C + A \cdot B \cdot \bar{C}$$

$$\begin{aligned} \bar{Y} &= \overline{A \cdot \bar{B} \cdot C + A \cdot B \cdot \bar{C}} \\ &= \overline{(A \cdot \bar{B} \cdot C)} \cdot \overline{(A \cdot B \cdot \bar{C})} \\ &= (\bar{A} + B + \bar{C}) \cdot (\bar{A} + \bar{B} + C) \end{aligned}$$

called a "product of sums"

How to write  $Y$  as a "product of sums" ?

First, write its complement  $\overline{Y}$  as a sum of products.

Because of time constraints, I decided to skip this example in the lecture.

You should go over it on your own.

A	B	C	Y	$\overline{Y}$
0	0	0	0	1
0	0	1	0	1
0	1	0	0	1
0	1	1	0	1
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	0	1

$$\overline{Y} = \overline{A} \cdot \overline{B} \cdot \overline{C} + \overline{A} \cdot \overline{B} \cdot C + \overline{A} \cdot B \cdot \overline{C} + \overline{A} \cdot B \cdot C + A \cdot \overline{B} \cdot \overline{C} + A \cdot B \cdot C$$

Then write  $Y = \overline{\overline{Y}}$  and apply de Morgan's Law.

$$\overline{Y} = \overline{A} \cdot \overline{B} \cdot \overline{C} + \overline{A} \cdot \overline{B} \cdot C + \overline{A} \cdot B \cdot \overline{C} + \overline{A} \cdot B \cdot C + A \cdot \overline{B} \cdot \overline{C} + A \cdot B \cdot C$$

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$$\overline{\overline{Y}} = (\overline{A} \cdot \overline{B} \cdot \overline{C} + \overline{A} \cdot \overline{B} \cdot C + \overline{A} \cdot B \cdot \overline{C} + \overline{A} \cdot B \cdot C + A \cdot \overline{B} \cdot \overline{C} + A \cdot B \cdot C)$$

$$= (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C})$$

$$= (A + B + C) \cdot (A + B + \overline{C}) \cdot (A + \overline{B} + C) \cdot (A + \overline{B} + \overline{C}) \cdot (\overline{A} + B + C) \cdot (\overline{A} + \overline{B} + \overline{C})$$

Sometimes we have expressions where various combinations of input variables give the same output. In the example below, if A is false then any combination of B and C will give the same output (namely true).

$$Y = A \cdot \bar{B} \cdot C + \bar{A}$$

A	B	C	$A \cdot \bar{B} \cdot C$	$\bar{A}$	Y
0	0	0	0	1	1
0	0	1	0	1	1
0	1	0	0	1	1
0	1	1	0	1	1
1	0	0	0	0	0
1	0	1	0	0	0
1	1	0	0	0	0
1	1	1	0	0	0

# Don't Care

We can simplify the truth table in such situations.

$$Y = A \cdot \bar{B} \cdot C + \bar{A}$$

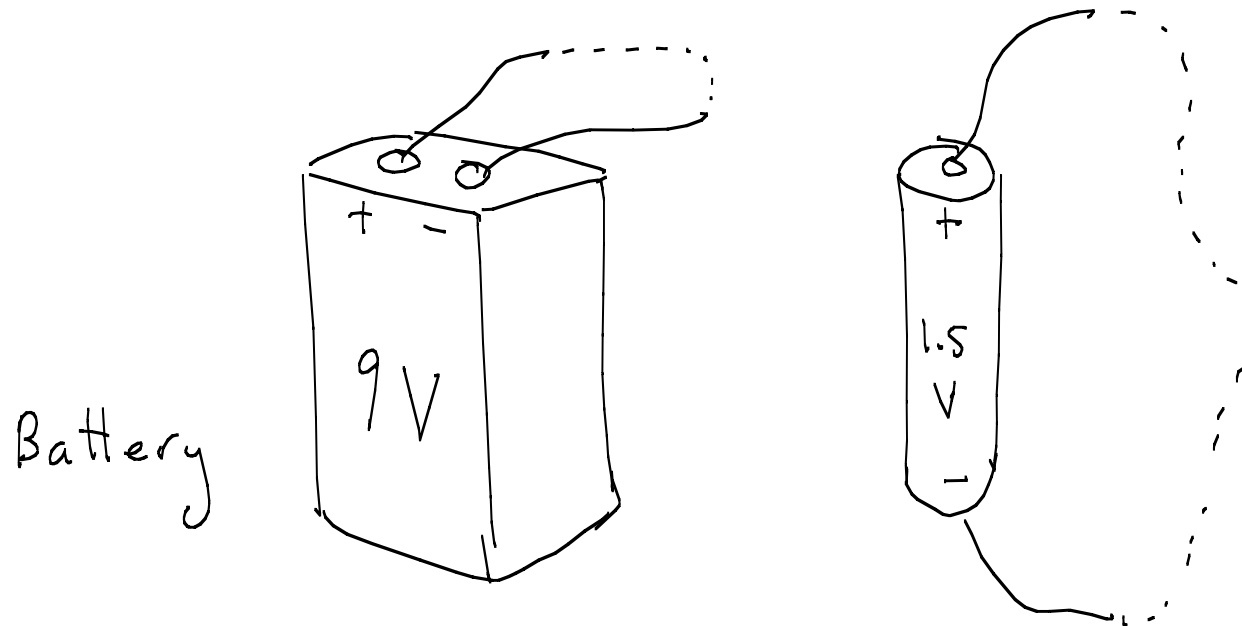
A	B	C	Y
0	X	X	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	0

~~X~~ means we "don't care" what values are there.



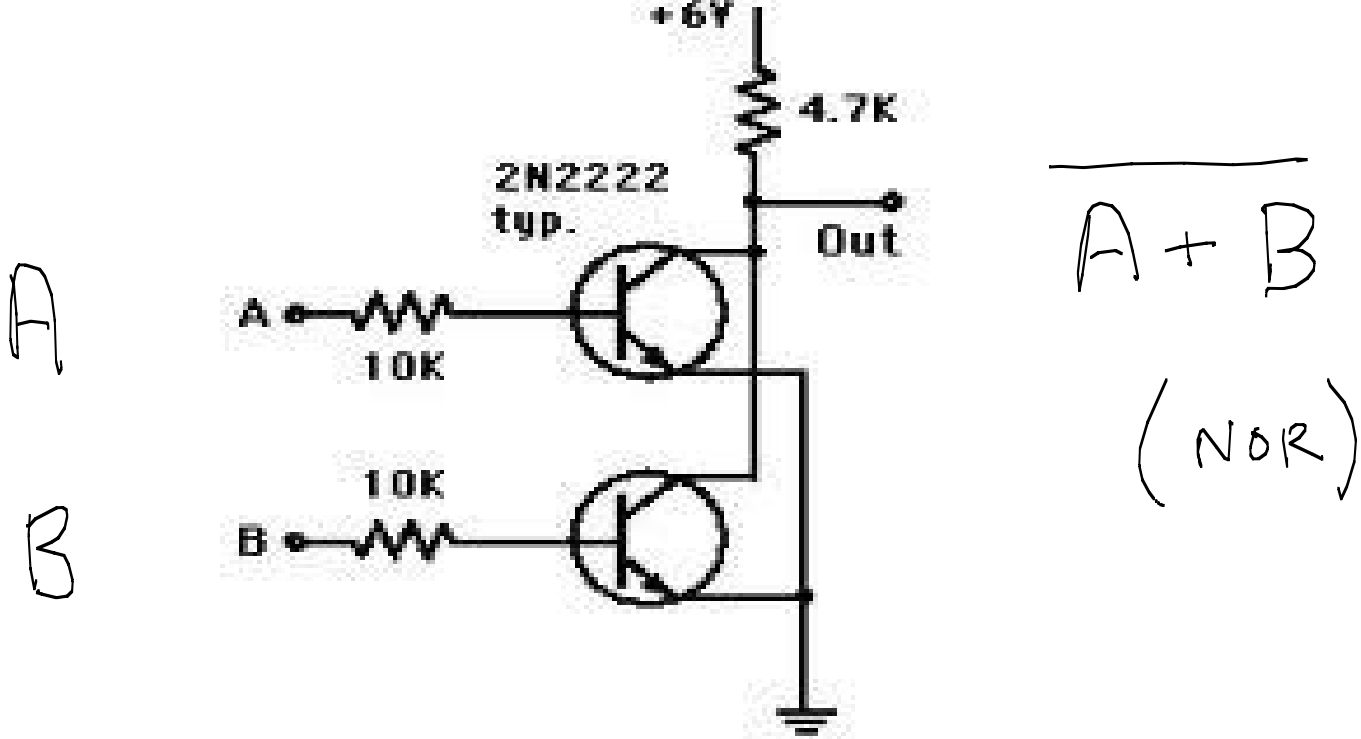
# What are the 0's and 1's in a computer?

A wire can have a voltage difference between two terminals, which drives current.



In a computer, wires can have two voltages:  
high (1, current ON) or low (0, current ~OFF)

Using circuit elements called "transistors" and "resistors", one can built circuits called "gates" that compute logical operations.



For each of the OR, AND, NAND, XOR gates, you would have a different circuit.

# Moore's Law (Gordon Moore was founder of Intel)

The number of transistors per  $\text{mm}^2$  approximately doubles every two years. (1965)

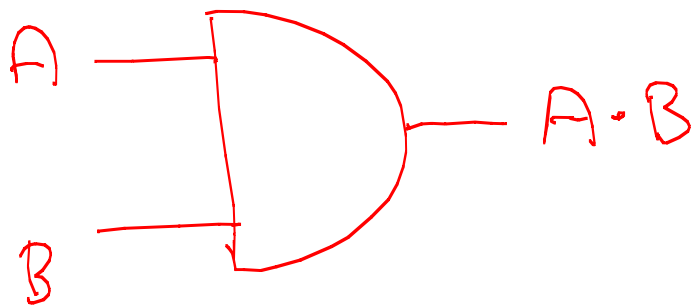
*It is an observation, not a physical law.*

It still holds true today, although people think that this cannot continue, because of limits on the size of atom and laws of quantum physics.

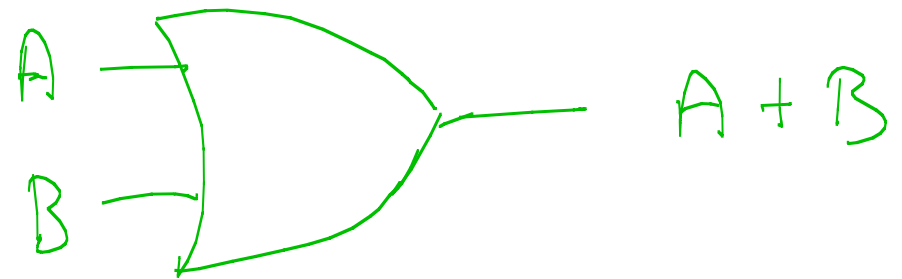
<http://phys.org/news/2015-07-law-years.html>

# Logic Gates

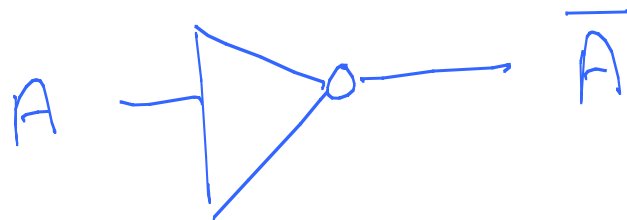
AND



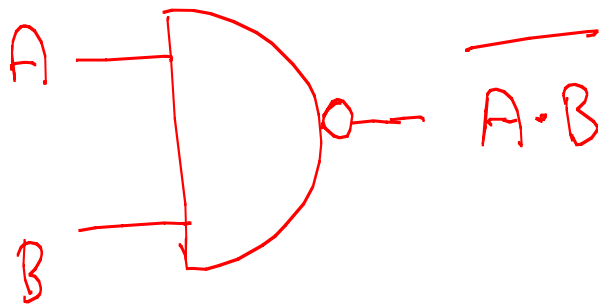
OR



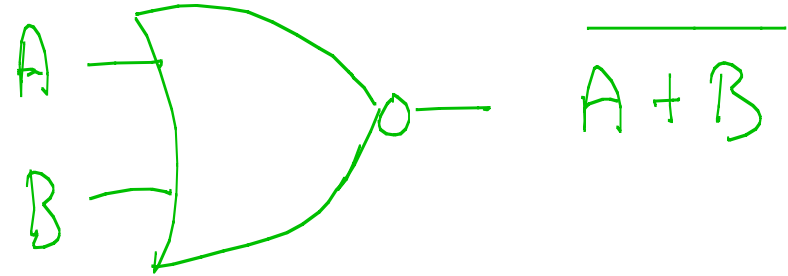
NOT



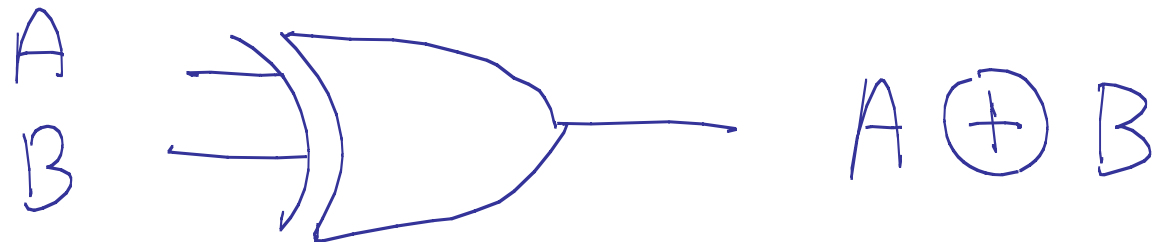
NAND



NOR



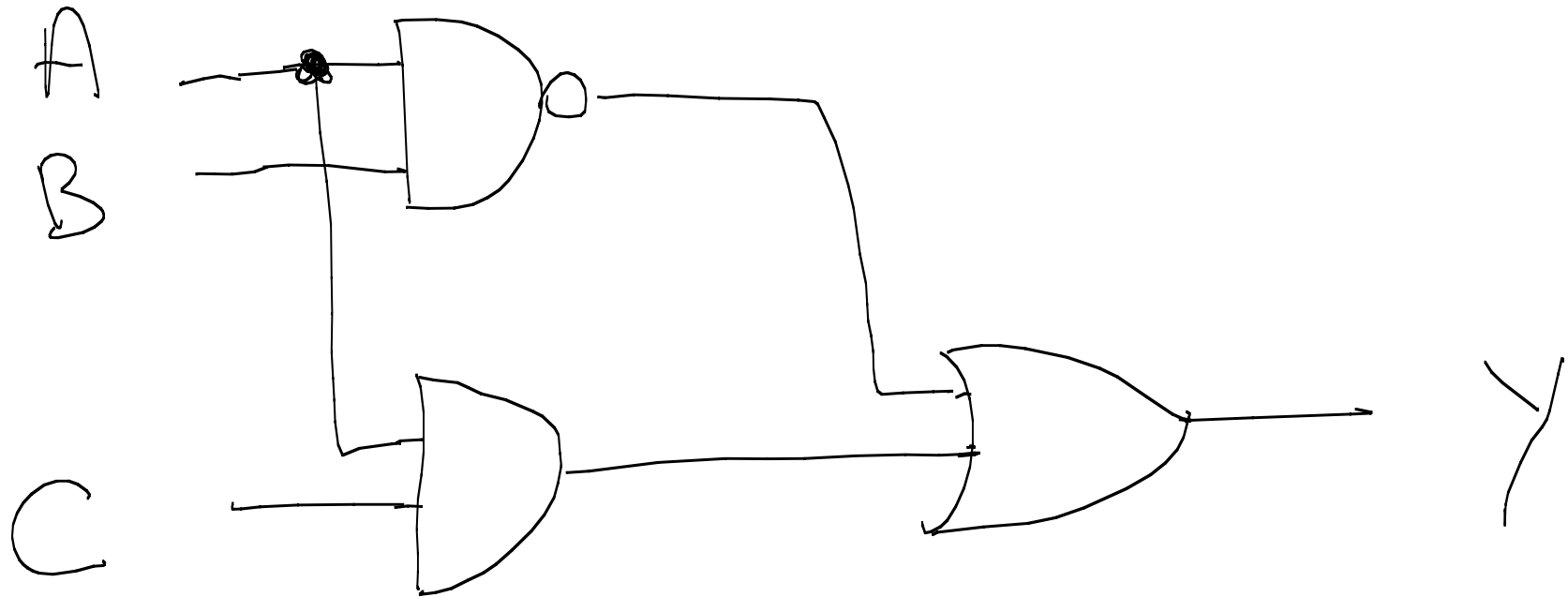
XOR



# Logic Circuit

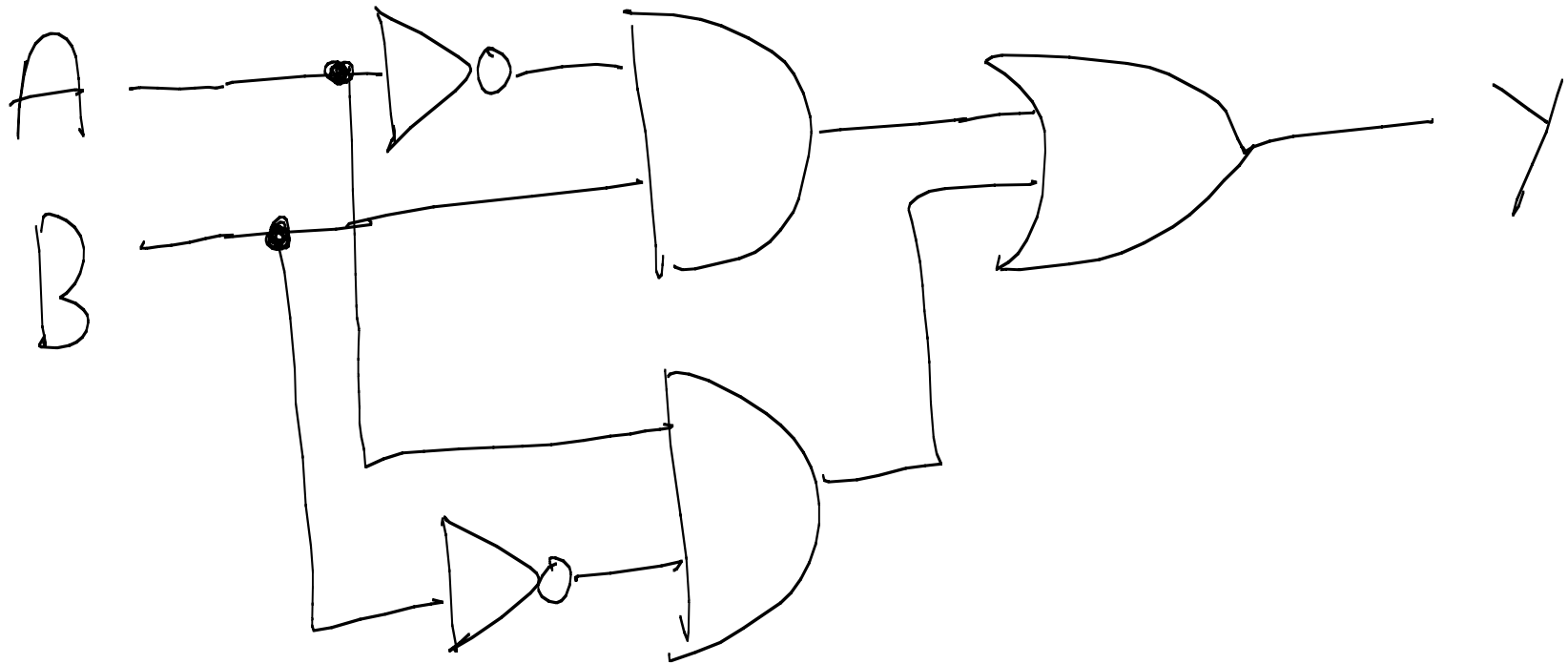
Example:

$$Y = \overline{A \cdot B} + A \cdot C$$



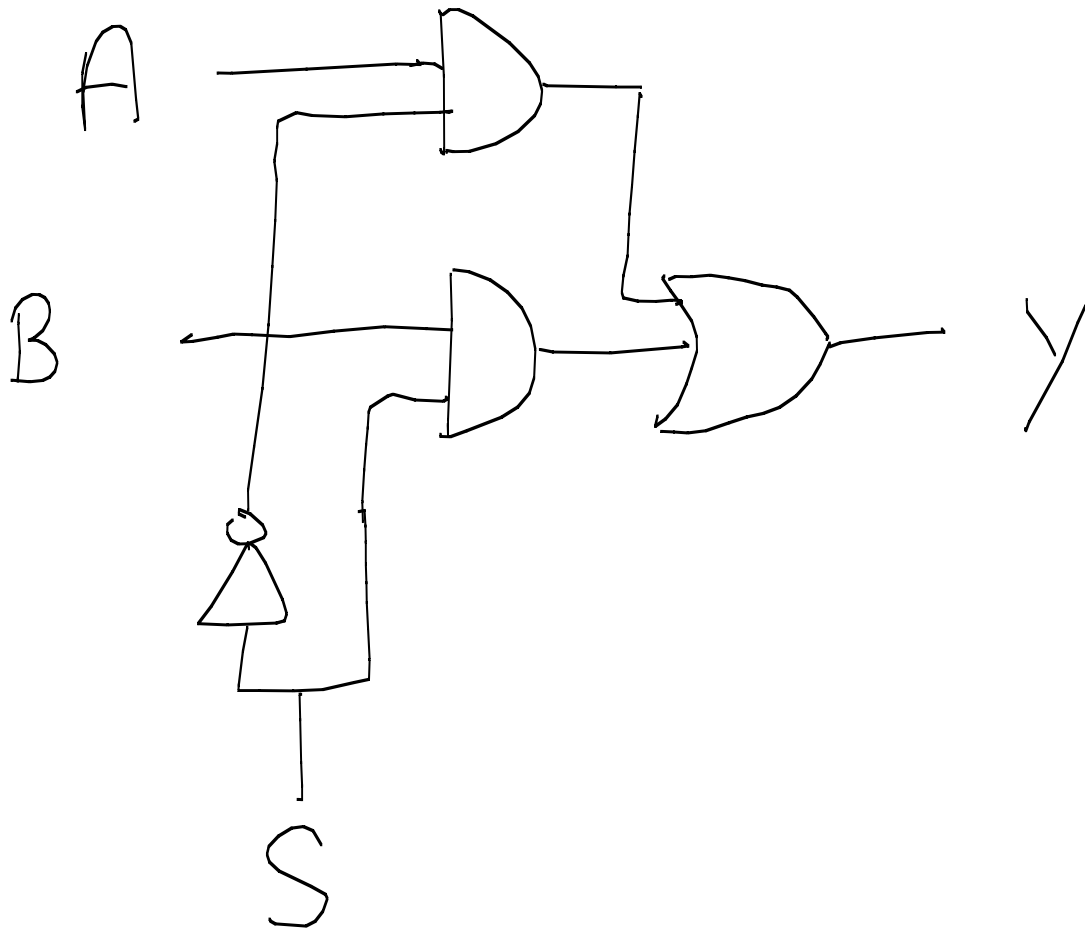
Example: XOR without using an XOR gate

$$Y = \bar{A} \cdot B + A \cdot \bar{B} = A \oplus B$$



# Multiplexor (selector)

$$Y = \overline{S} \cdot A + S \cdot B$$

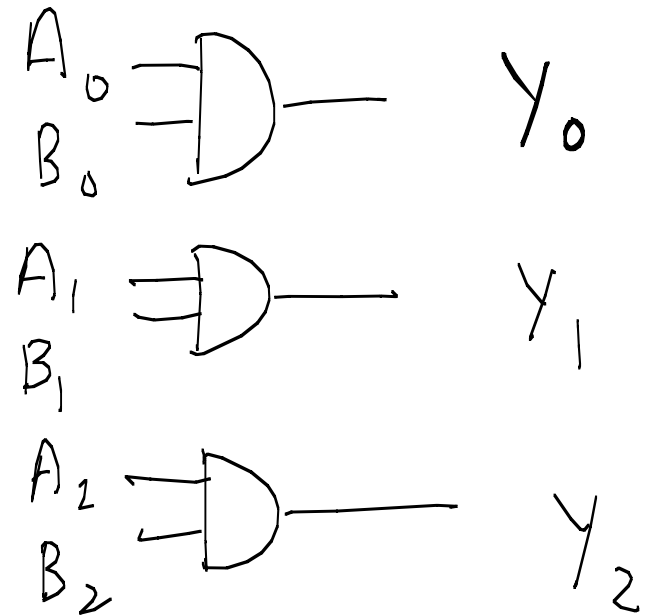
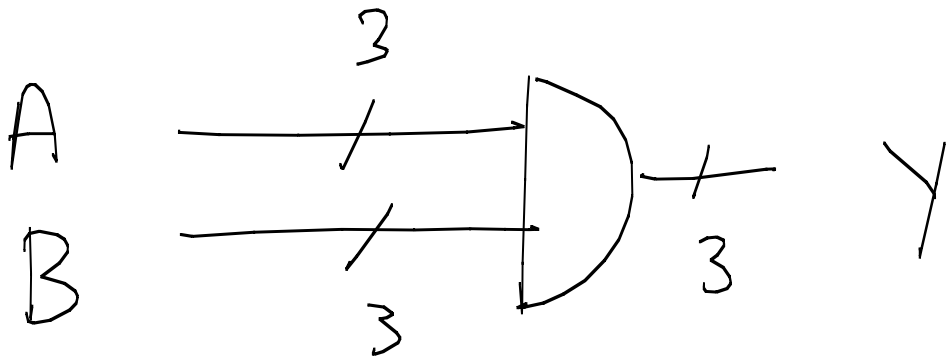


if S  
    Y = B  
else  
    Y = A



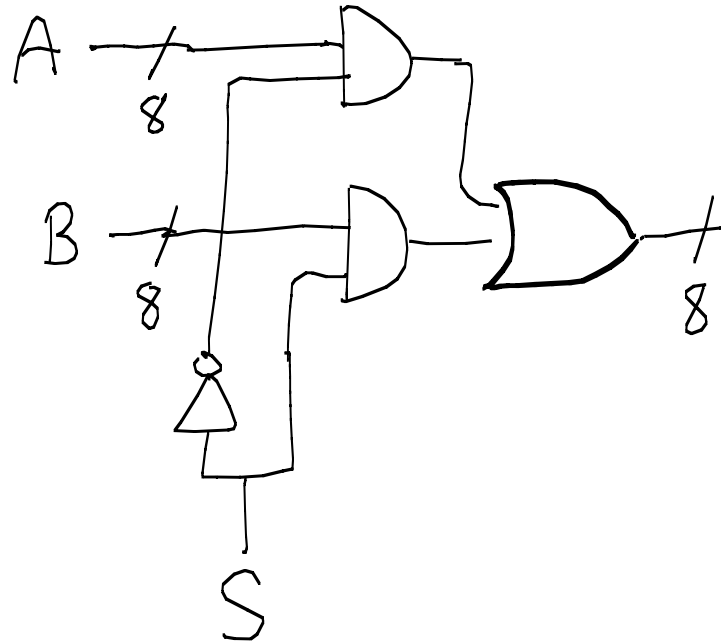
# Notation

Suppose  $A$  and  $B$  are each 3 bits ( $A_2 A_1 A_0$ ,  $B_2 B_1 B_0$ )



Suppose A and B are each 8 bits ( $A_7 A_6 \dots A_0$ ,  $B_7 B_6 \dots B_0$ )  
We can define an 8 bit multiplexor (selector).

Notation:



In fact we would build this from 8 separate one-bit multiplexors.

Note that the selector S is a single bit. We are selecting either all the A bits or all the B bits.

# Announcement

The enrollment cap will be lifted before DROP/ADD to allow students on the waitlist to register.